Toward a Salmon Conjecture

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...and why even a vegetarian might care...

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The salmon prize

In 2007, E. Allman offered a prize of Alaskan salmon (!) to whoever finds the defining ideal of $\sigma_4(\mathbb{P}^3 \times \mathbb{P}^3 \times \mathbb{P}^3)$.

This algebraic variety may be viewed as a statistical model for evolution.



Nucleotides $\{A, C, G, T\}$ Independent extant species Unknown (hidden) Ancestor Invariants of this statistical model

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 $\Delta \subset \mathbb{P}\mathbb{R}^{64} \subset \mathbb{P}\mathbb{C}^{64}$

Nucleotides $\{A, C, G, T\} \leftrightarrow \mathbb{PC}^4$. Independent extant species $\leftrightarrow Seg(\mathbb{P}^3 \times \mathbb{P}^3 \times \mathbb{P}^3)$. Unknown (hidden) Ancestor $\leftrightarrow 4^{th}$ secant variety. Invariants of this statistical model \leftrightarrow ideal of the algebraic variety.

The salmon prize

Allman's Motivation: Work of Allman-Rhodes'03 implies that solving the salmon problem would provide all phylogenetic invariants for a whole class of binary evolutionary tree models!

As in this example, nice varieties in spaces of tensors (like secant varieties) appear in several fields outside of mathematics, such as

- algebraic statistics (other problems like this one)
- computational complexity theory (bounding the complexity of algorithms via ranks of tensors)
- signal processing (CDMA protocol for mobile phones)
- physics (quantum information theory and measures of entanglement)
- computer vision (multi-view geometry)
- ... your favorite variety?

Recent history and current status

- [Landsberg–Manivel 2004]: Some equations of $\sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3)$ in degrees 5,6,9 in representation theoretic language.
- [Landsberg-Manivel 2008]: Reduced set-theoretic problem for $\sigma_4(\mathbb{P}^3 \times \mathbb{P}^n \times \mathbb{P}^m)$, $n, m \ge 3$ to $\sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3)$.
- <u>October 2008</u> Sturmfels asked for explicit (M2) version of degree 6 equations. (Now an ancillary file on ArXiv [Bates–O. 2011]).
- <u>December 2008</u> (O–) MSRI Algebraic Statistics Workshop: Conjecture about zero set of degree 6 equations, if confirmed would prove set-theoretic result.
- March 2010 (Friedland): set-theoretic result using degrees 5, 9, 16. Second version corrects proof of Landsberg-Manivel reduction.
- July 2010 (Bates): Numerical Algebraic Geometry (NAG) calculation in Bertini for deg. 6 equations, MSRI conjecture ⇒ numerical, set-theoretic result using degrees 5, 6, 9, [Bates-O 2011].
- April 2011 (Friedland–Gross): Explicit equations + previous proof of Friedland: confirm NAG result without numerical methods.
- Ideal theoretic problem is still open.

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Secant varieties

Let $A = \{a_i\}, B = \{b_j\}, C = \{c_k\}$, be \mathbb{C} -vector spaces, then the tensor product is $A \otimes B \otimes C = \{a_i \otimes b_j \otimes c_k\}$, with coordinates p_{ijk} .

• Segre variety (rank 1 tensors): (Independence model) Defined by

$$\begin{aligned} & \textit{Seg}: \mathbb{P}A \times \mathbb{P}B \times \mathbb{P}C & \longrightarrow & \mathbb{P}\big(A \otimes B \otimes C\big) \\ & & ([a], [b], [c]) & \longmapsto & [a \otimes b \otimes c]. \end{aligned}$$

• The r^{th} secant variety of a variety $X \subset \mathbb{P}^n$: (Mixture model)

$$\sigma_r(X) = \bigcup_{x_1, \dots, x_r \in X} \mathbb{P}(\operatorname{span}\{x_1, \dots, x_r\}) \subset \mathbb{P}^n.$$

A useful reduction

Theorem (Landsberg-Manivel '08, Friedland'10)

 $\sigma_4(\mathbb{P}^3 \times \mathbb{P}^3 \times \mathbb{P}^3)$ is the zero set of:

- $M_5 = \{ (Strassen's [1983] degree 5 commutation conditions) \}$
- 2 Equations inherited from $\sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3)$
 - Key point: It remains to find the equations of $\sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3)!$
 - Note: M_5 is a 1728 dimensional irreducible *G*-module, for $G = GL(4) \times GL(4) \times GL(4) \rtimes \mathfrak{S}_3$ with a natural basis of polynomials with 180 or 360 or 540 monomials (see also [Allman-Rhodes '03]).

Symmetry

• The symmetry group of the salmon variety

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\sigma_4(\mathbb{P}A\times\mathbb{P}B\times\mathbb{P}C)
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is change of coordinates in each factor,

 $GL(A) \times GL(B) \times GL(C)$

(or $GL(A) \times GL(B) \times GL(C) \rtimes \mathfrak{S}_3$ when $A \cong B \cong C$).

- Good news: A large group acts and we can use tools from Representation Theory!
- This symmetry is a powerful tool and we should exploit it!

Representation Theory notation

- Module notation: $S^d(A \otimes B \otimes C) = \mathbb{C}[p_{ijk}]_d$.
- Fact: $S^d(A \otimes B \otimes C)$ is a $GL(A) \times GL(B) \times GL(C)$ -module.
- The irreducible submodules of S^d(A ⊗ B ⊗ C) are isomorphic to Schur modules indexed by certain partitions π₁, π₂, π₃ of d:

$$S_{\pi_1}A\otimes S_{\pi_2}B\otimes S_{\pi_3}C,$$

and usually occur with multiplicity - this makes us work harder.

• Given π_1, π_2, π_3 and the multiplicity, there is a combinatorial algorithm for constructing polynomials!

An ideal membership test

Apply [Landsberg–Manivel'04] ideal membership test: For each d,

- decompose $S^d(A^* \otimes B^* \otimes C^*)$ as a $GL(A) \times GL(B) \times GL(C)$ -module.
- for each module (isotypic component), test a highest weight vector (highest weight space) for vanishing on σ₄(ℙA × ℙB × ℙC).
- output: $\mathcal{I}_d(\sigma_4(\mathbb{P}A \times \mathbb{P}B \times \mathbb{P}C))$ as a list of modules.

Works well for small degree and produced the following results:

•
$$\mathcal{I}_5(\sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3)) = 0.$$

•
$$\mathcal{I}_6(\sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3)) = S_{2,2,2}\mathbb{C}^3 \otimes S_{2,2,2}\mathbb{C}^3 \otimes S_{3,1,1,1}\mathbb{C}^4$$

= { ten degree 6 polynomials on 36 variables }.

Inheritance via an example

Proposition (example of Proposition 4.4 Landsberg–Manivel'04) $\widetilde{M}_{6} := S_{(2,2,2)} \mathbb{C}^{4} \otimes S_{(2,2,2)} \mathbb{C}^{4} \otimes S_{(3,1,1,1)} \mathbb{C}^{4} \subset \mathcal{I}(\sigma_{4}(\mathbb{P}^{3} \times \mathbb{P}^{3} \times \mathbb{P}^{3}))$ *if and only if* $M_{6} := S_{(2,2,2)} \mathbb{C}^{3} \otimes S_{(2,2,2)} \mathbb{C}^{3} \otimes S_{(3,1,1,1)} \mathbb{C}^{4} \subset \mathcal{I}(\sigma_{4}(\mathbb{P}^{2} \times \mathbb{P}^{2} \times \mathbb{P}^{3})).$

Note: dim $(\tilde{M}_6) = 10^3$ but dim $(M_6) = 10$, and has basis of polynomials, each with 576 or 936 monomials.

At every stage we study the smallest module possible. This is a significant dimension reduction.

For $\sigma_4(\mathbb{P}^3 \times \mathbb{P}^3 \times \mathbb{P}^3)$ we only need to consider $S_{\pi_1}A \otimes S_{\pi_2}B \otimes S_{\pi_3}C$ where π_1, π_2, π_3 have 4 parts, and those equations we get from inheritance.

What is a flattening?

Express a tensor $T = \sum_{i,j,k} p_{ijk} a_i \otimes b_j \otimes c_k \in A \otimes B \otimes C$ as a matrix: $T = \sum_{i} a_i \otimes \left(\sum_{i,k} p_{ijk} b_j \otimes c_k \right) \in A \otimes (B \otimes C)$ For example: $\mathbb{C}^3 \otimes \mathbb{C}^3 \otimes \mathbb{C}^3 \cong \mathbb{C}^3 \otimes (\mathbb{C}^3 \otimes \mathbb{C}^3) \cong \mathbb{C}^3 \otimes \mathbb{C}^9$: $T = [p_{ijk}] = \sum_i a_i \otimes (\sum_{i,k} p_{ijk} b_j \otimes c_k) = \sum_i a_i \otimes X_i$ $\begin{array}{ccc} p_{113} & p_{123} & p_{133} \\ p_{213} & p_{223} & p_{233} \end{array} \right)$ $\psi_{0,T} = \begin{pmatrix} p_{111} & p_{121} & p_{131} & | & p_{112} & p_{122} \\ p_{211} & p_{221} & p_{231} & | & p_{212} & p_{222} \\ p_{311} & p_{321} & p_{331} & | & p_{312} & p_{322} \end{pmatrix}$ p_{132}

 p_{232} *p*₃₁₂ P332 *p*₃₁₃ p₃₂₃ P333

 $= (X_1 \mid X_2 \mid X_3)$

When they exist, $(r+1) \times (r+1)$ minors of $\psi_{0,T}$ are (some) equations of $\sigma_r(\mathbb{P}A \times \mathbb{P}B \times \mathbb{P}C).$

Flattenings and subspace varieties

Tensors that can be written using fewer variables:

$$Sub_{p,q,r} := \left\{ [T] \in \mathbb{P}(A \otimes B \otimes C) \mid \begin{array}{c} \exists \mathbb{C}^p \subseteq A, \exists \mathbb{C}^q \subseteq B, \exists \mathbb{C}^r \subseteq C, \\ \text{and } [T] \in \mathbb{P}(\mathbb{C}^p \otimes \mathbb{C}^q \otimes \mathbb{C}^r) \end{array} \right\}$$

Theorem (3.1, Landsberg–Weyman '07)

 $Sub_{p,q,r}$ is normal with rational singularities. Its ideal is generated by the minors of flattenings;

$$\left(\bigwedge^{p+1}A\otimes \bigwedge^{p+1}(B\otimes C)\right)\oplus \left(\bigwedge^{q+1}B\otimes \bigwedge^{q+1}(A\otimes C)\right) \oplus \left(\bigwedge^{r+1}(A\otimes B)\otimes \bigwedge^{r+1}C\right)$$

Fact: $Sub_{r,r,r} \supseteq \sigma_r(\mathbb{P}A \times \mathbb{P}B \times \mathbb{P}C)$

Key Point: The subspace varieties contain secant varieties, and therefore they give some of the equations of the secant varieties.

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A result of Strassen

Theorem (Strassen 1988 (reinterpreted by Landsberg–Manivel)) The ideal of the hypersurface $\sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^2) \subset \mathbb{P}^{26}$ is generated in degree 9 by a nonzero vector in the 1 dimensional module

 $\mathcal{S}_{(3,3,3)}\mathbb{C}^3\otimes\mathcal{S}_{(3,3,3)}\mathbb{C}^3\otimes\mathcal{S}_{(3,3,3)}\mathbb{C}^3$

Since $\sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^2) \subset \sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3)$, inheritance implies that $M_9 := S_{(3,3,3)} \mathbb{C}^3 \otimes S_{(3,3,3)} \mathbb{C}^3 \otimes S_{(3,3,3)} \mathbb{C}^4 \subset \mathcal{I}(\sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3))$

Strassen's polynomial *only* has 9,216 monomials on 27 variables. $\dim(M_9) = 20$, with natural basis of polynomials with 9,216 or 25,488 or 43,668 monomials on 36 variables! 23 Mb file of polynomials... :-(Strassen's equation: A useful reformulation by Ottaviani

 $T = [p_{ijk}] = \sum_i a_i \otimes (\sum_{j,k} p_{ijk} b_j \otimes c_k) = \sum_i a_i \otimes X_i$ Strassen's equation is the determinant of the 9 × 9 matrix:

$$\psi_{\mathcal{T}} = \left(\begin{array}{ccc} 0 & X_3 & -X_2 \\ -X_3 & 0 & X_1 \\ X_2 & -X_1 & 0 \end{array}\right)$$

Basic idea:

$$\psi_{1,T+T'} = \psi_T + \psi_{T'} \qquad \text{construction is linear in } T$$

$$Rank(T) = 1 \Rightarrow Rank(\psi_T) = 2 \qquad \text{base case}$$

$$\therefore Rank(T) = r \Rightarrow Rank(\psi_T) \le 2r \qquad \text{upper bound on rank}$$

Numerical Algebraic Geometry: Bertini

Theorem*

The zero set of M_6 (ten polynomials on 36 variables) has precisely two components of codimensions 4 and 6 and degrees 345 and 84 respectively. $\mathcal{V}(M_6) = \sigma_4 (\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3) \cup Sub_{3,3,3}.$

Using numerical homotopy continuation (Bertini):

- Original computation (July 2010) 2 weeks of computational time on 8 processors: 2.66GHz × 8p × 336h = 7150GHzh
- Regeneration -Hauenstein, Sommese & Wampler, May 2011. $2.33GHz \times 65p \times 20h = 3029GHzh - confirmed same result$
- Small tracking and final tolerances $(10^{-10} \text{ or smaller})$
- Adaptive precision numerical methods
- Checks and error controls built into Bertini such as checking at t = 0.1 that no paths have crossed.

Confirms conjecture from MSRI 2008.

Numerical Algebraic Geometry: Bertini

Theorem*

 $\mathcal{V}(M_6 + M_9) = \sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3).$

Suppose $x \in \mathcal{V}(M_6) = \sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3) \cup Sub_{3,3,3}$. If $x \notin \sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3)$, then use M_9 and consider $x \in Sub_{3,3,3} \cap \mathcal{V}(M_9)$ $\Rightarrow x$ is in some $\sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^2) \subset \sigma_4(\mathbb{P}^2 \times \mathbb{P}^2 \times \mathbb{P}^3)$.

Theorem* (Corollary to Landsberg-Manivel 2008, Friedland 2010)

The salmon variety is cut out set-theoretically in degrees 5, 6, 9:

$$\mathcal{V}\left(\mathit{M}_{5}+\tilde{\mathit{M}}_{6}+\tilde{\mathit{M}}_{9}
ight)=\sigma_{4}\left(\mathbb{P}^{3} imes\mathbb{P}^{3} imes\mathbb{P}^{3}
ight)$$

Resolves the salmon problem set-theoretically. Provides a more efficient set of equations than [Friedland 2010]. Sharpens the conjecture for the ideal-theoretic question. *Friedland–Gross 2011 make Theorem* into Theorem.*

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A template for finding equations of varieties coming from applications

The salmon variety has been studied via the following:

- Start: statistical model, space of special tensors, etc.
- **2** Find the corresponding algebraic variety X.
- Solution Find the largest symmetry group G acting X.
- Study $\mathcal{I}(X)$ as a *G*-module using Representation Theory.
- Sompute all modules in $\mathcal{I}_d(X)$ for small degree (nec. conditions).
- **1** Use Numerical Algebraic Geometry to compute unknown zero-sets.
- Try to make geometric reductions to show that the known invariants suffice.
- Try to prove what you know* is true.